

# The internal language of sheaves

## Applications to algebraic geometry

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# Outline

- 1 Internal Language
- 2 Kripke-Joyal Semantics
- 3 Modal Operators
- 4 Relation between formulas on  $X_{\square}$  and formulas on  $X +$  why it is unclear
- 5 Solution

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# Internal language of a topos

A topos is a category that has all the properties of **Set** required for finitary reasoning. That is, we can do first order logic inside a topos.

## Definition (Elementary topos)

An elementary topos is a category that

- has all finite limits
- is cartesian closed
- has a subobject classifier  $\Omega$

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## Example

The category of sheaves over a topological space  $X$ ,  $\text{Sh}(X)$ , is a topos.

# Internal language, context

To be able to do first order logic, we must be able to represent judgment  
i.e.  $\Gamma \vdash F$ .

## Context

- Objects  $\mathcal{F}$  are our types
- Variables  $x$  of type  $\mathcal{F}$  are interpreted as the identity morphism  $\text{id} : \mathcal{F} \rightarrow \mathcal{F}$
- Terms  $\sigma$  of type  $\mathcal{F}$  are interpreted as morphisms  $\sigma = (\sigma_1, \dots, \sigma_n) : \mathcal{E}_1 \times \dots \times \mathcal{E}_n \rightarrow \mathcal{F}$ .
- Formulae are the terms  $\mathcal{F} \rightarrow \Omega$

So  $\Gamma \vdash F$ , will be represented by  $\varphi : \mathcal{E}_1 \times \dots \times \mathcal{E}_n \rightarrow \Omega$ .

# Internal language

Moreover, we need to build constructors corresponding to  $\perp$ ,  $\top$ ,  $\wedge$ ,  $\vee$ ,  $\Rightarrow$ ,  $\forall$  and  $\exists$ .

## Example (Connectives)

For any formulas  $\varphi : \mathcal{F} \rightarrow \Omega$  and  $\psi : \mathcal{F} \rightarrow \Omega$  we want to build a new formula

$$\begin{aligned}\varphi \wedge \psi &: \mathcal{F} \xrightarrow{\langle \varphi, \psi \rangle} \Omega \times \Omega \xrightarrow{\wedge_{\Omega}} \Omega \\ \varphi \vee \psi &: \mathcal{F} \xrightarrow{\langle \varphi, \psi \rangle} \Omega \times \Omega \xrightarrow{\vee_{\Omega}} \Omega \\ \varphi \Rightarrow \psi &: \mathcal{F} \xrightarrow{\langle \varphi, \psi \rangle} \Omega \times \Omega \xrightarrow{\Rightarrow_{\Omega}} \Omega\end{aligned}$$

## Proposition

*We can also define  $\top$ ,  $\perp$ ,  $\exists$ ,  $\forall$*

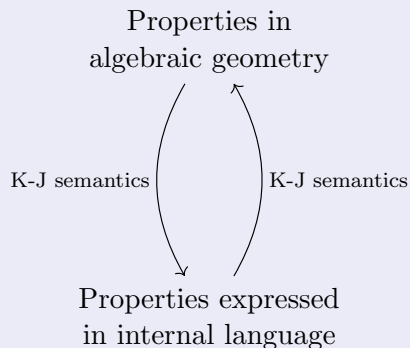
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# Kripke-Joyal semantics

## Property transfer idea



Properties we want  
to prove  
**Complicated objects**

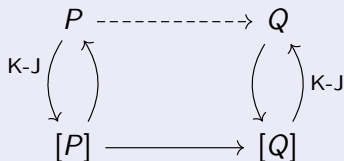
Translated  
properties  
**Simpler objects**

# Internal proofs 101

## General scheme of internal proofs

Suppose we have a properties  $P$  and  $Q$  and we want to give a proof of  $P \rightarrow Q$ . Then we can translate  $P$  and  $Q$  to their corresponding internal statements  $[P]$  and  $[Q]$  and give a proof of  $[P] \rightarrow [Q]$  instead.

External



Internal

# Kripke-Joyal semantics

## Notation

*Given a formula  $\varphi : \mathcal{F} \rightarrow \Omega$  and an open set  $U \subseteq X$ , we have  $\varphi_U : \mathcal{F}(U) \rightarrow \Omega(U)$ . Given  $\alpha \in \mathcal{F}(U)$ , we will write  $\varphi(\alpha)$  instead of  $\varphi_U(\alpha)$*

## Notation

*For a formula  $\varphi : \mathcal{F} \rightarrow \Omega$  then we will use  $\{x \mid \varphi(x)\}$  to denote the sheaf it classifies*

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## Definition (Kripke-Joyal semantics)

Let  $X$  be a topological space,  $U \subseteq X$  and  $\varphi$  be some formula  $\varphi : \mathcal{F} \rightarrow \Omega$  in the internal language of  $\text{Sh}(X)$  then.

$$U \models \varphi(\alpha) \text{ for } \alpha \in \mathcal{F}(U) \quad \text{iff} \quad \alpha \in \{x | \varphi(x)\}(U)$$

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## Example (Recursive unwinding)

- $U \vDash (\varphi \wedge \psi)(\alpha)$  iff  $U \vDash \varphi(\alpha)$  and  $U \vDash \psi(\alpha)$ .
- $U \vDash (\varphi \Rightarrow \psi)(\alpha)$  iff for all open  $V \subseteq U$   $V \vDash \varphi(\alpha|_V)$  implies  $V \vDash \psi(\alpha|_V)$ .

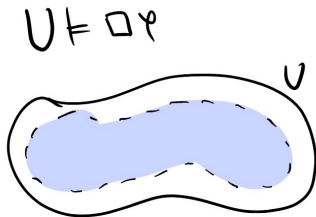
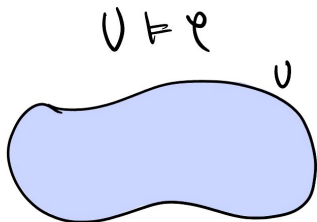
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# Modalities

Sometimes having a property holding globally is too strict, we also want to be able to describe "local" properties.

We want some operator  $\square : \Omega \rightarrow \Omega$  that does this:





# Modal operators

## Definition (Modal operator)

A modal operator is a map  $\square : \Omega \rightarrow \Omega$  such that

- 1  $\varphi \implies \square\varphi$
- 2  $\square\square\varphi \implies \square\varphi$
- 3  $\square(\varphi \wedge \psi) \iff \square\varphi \wedge \square\psi$

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## Examples (Some general modal operators)

For a fixed formula in the internal language  $\alpha$ , the following are modal operators

- 1  $\square\varphi := (\alpha \Rightarrow \varphi)$
- 2  $\square\varphi := ((\varphi \Rightarrow \alpha) \Rightarrow \alpha)$
- 3  $\square\varphi := \neg\neg\varphi$

## Specific modal operators and their translations

Since the open set  $U \subseteq X$  is an element of  $\Omega(X)$  we have that  $V \vDash U \iff V \subseteq U$ .

### Definition ( $!x$ )

Because  $x \notin V \iff V \vDash \text{int}(X \setminus \{x\})$  we let  $!x$  denote  $\text{int}(X \setminus \{x\})$ .

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## Proposition (Instances of modal operators)

For  $x \in X$  we have that for all open  $V \subseteq X$

- $V \vDash \neg\neg\varphi \iff \exists W \subseteq V$  which is open and dense such that  $W \vDash \varphi$
- $V \vDash ((\varphi \Rightarrow !x) \Rightarrow !x) \iff x \notin V$  or  $\exists W \subseteq V$  with  $x \in W$  and  $W \vDash \varphi$

## $X_{\square}$ , the associated map $j$

Sometimes studying these local properties directly on  $X$  can be hard, thus we will define a subspace,  $X_{\square}$ , which validates the properties on which they are easier to study.

## $X_{\Box}$ , the associated map $j$

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### Definition (Associated map to a modal operator)

A modal operator induces a map on the elements of  $\Omega(X)$  (i.e. the open subsets of  $X$ )

$$\begin{aligned} j : \Omega(X) &\longrightarrow \Omega(X) \\ U &\longmapsto \bigcup \{V \subseteq X \mid V \text{ open, } V \models \Box U\} \end{aligned}$$

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### Proposition (Properties of $j$ )

*In particular we have that for any open  $U, V \subseteq X$*

- 1  $U \subseteq j(U)$
- 2  $j(j(U)) \subseteq j(U)$
- 3  $j(U \cap V) = j(U) \cap j(V)$

## $X_{\square}$ as an induced subspace

### Definition (Subspace associated to a modal operator)

$j$  defines a subspace of  $X$ , denoted by  $X_{\square}$ . It has a frame of opens  
 $\mathcal{T}(X_{\square}) := \{U \in \Omega(X) \mid j(U) = U\}$



## $X_{\square}$ as an induced subspace

### Definition (Subspace associated to a modal operator)

$j$  defines a subspace of  $X$ , denoted by  $X_{\square}$ . It has a frame of opens  $\mathcal{T}(X_{\square}) := \{U \in \Omega(X) \mid j(U) = U\}$

Modal operator	associated map	$X_{\square}$
$\square\varphi \equiv (U \Rightarrow \varphi)$	$j(V) = \text{int}(U^c \cup V)$	$U$
$\square\varphi \equiv \neg\neg\varphi$	$j(V) = \text{int}(\text{cl}(V))$	smallest dense sublocale of $X$
$\square\varphi \equiv ((\varphi \Rightarrow !x) \Rightarrow !x)$	$j(V) = \begin{cases} X & \text{if } x \in V \\ X \setminus \text{cl}(\{x\}) & \text{if } x \notin V \end{cases}$	$\{x\}$

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# Relation between formulas in $\text{Sh}(X_{\square})$ and those in $\text{Sh}(X)$

## Theorem ( $\square$ 'd formula relation)

Let  $X$  be a topological space and  $\square$  be a modal operator on  $\text{Sh}(X)$  then for any formula  $\varphi$

$$X \models_{\text{Sh}(X)} \varphi^{\square} \iff X_{\square} \models_{\text{Sh}(X_{\square})} \varphi$$

where all parameters on the right side are pulled back to  $X_{\square}$  along  $X_{\square} \hookrightarrow X$ .

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## Proposition ( $\square$ 'd formula relation for geometric formulas)

For geometric formulas and geometric implications  $\varphi$  we have

$$\square\varphi \iff \varphi^{\square}$$

$$X \models_{\text{Sh}(X)} \square\varphi \iff X_{\square} \models_{\text{Sh}(X_{\square})} \varphi$$

# Algebraic geometry

Proposition ( $\square$ 'd formula relation for geometric formulas)

*For geometric formulas and geometric implications  $\varphi$  we have*

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Proposition ( $\square$ 'd formula relation for geometric formulas)

For geometric formulas and geometric implications  $\varphi$  we have

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Example (Properties on the space vs at stalks)

Let  $x \in X$  and  $\square\varphi := ((\varphi \Rightarrow !x) \Rightarrow !x)$ , then by the theorem

$$X \models \varphi^{\square} \iff \varphi \text{ holds at the stalk at } x$$

From the proposition we get that a geometric implication holds if and only if it holds at every stalk.

# Problem

## Theorem

Let  $X$  be a topological space and  $\square$  be a modal operator on  $\text{Sh}(X)$  then for any formula  $\varphi$

$$X \models_{\text{Sh}(X)} \varphi^\square \iff X_\square \models_{\text{Sh}(X_\square)} \varphi$$

where all parameters on the right side are pulled back to  $X_\square$  along  $X_\square \hookrightarrow X$ .

## Question

Since the formula  $\varphi$  in  $\text{Sh}(X)$  can have a context that is not a  $\square$ -sheaf as well as domains of quantification's that are not  $\square$ -sheaves, does the equivalence

$$X \models_{\text{Sh}(X)} \varphi^\square \iff X_\square \models_{\text{Sh}(X_\square)} \varphi$$

make sense?

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# Solution

## Outline of solution

- 1 We are going to define a construction  $+$  such that for all sheaves  $\mathcal{F}$ ,  $\mathcal{F}^{++}$  is a  $\square$ -sheaf.
- 2 We are going to define an associated translation on formulas such that for a formula  $\varphi$ ,  $\varphi^\square \Leftrightarrow (\varphi^\square)^{++}$
- 3 This will enable us to prove

$$X \models_{\text{Sh}(X)} (\varphi^\square)^{++} \iff X_\square \models_{\text{Sh}(X_\square)} \varphi^{++}$$

- 4 And we can conclude

$$X \models_{\text{Sh}(X)} \varphi^\square \iff X_\square \models_{\text{Sh}(X_\square)} \varphi$$

# Sheaves on $X_{\square}$

What do sheaves on  $X_{\square}$  look like?

## Definition ( $\square$ -sheaf)

A  $\square$ -sheaf is a sheaf in  $\text{Sh}(X)$  that "looks" like a sheaf on  $X_{\square}$

## Proposition

*There is an equivalence of categories  $\text{Sh}(X_{\square}) \simeq \text{Sh}_{\square}(\text{Sh}(X))$ , where  $\text{Sh}_{\square}(\text{Sh}(X))$  is the category of  $\square$ -sheaves in  $\text{Sh}(X)$ . It is induced by the following map*

$$\begin{aligned} i : \mathcal{O}(X) &\rightarrow \mathcal{T}(X_{\square}) \\ U &\mapsto j(U) \end{aligned}$$

# How to construct $\mathcal{F}^{++}$

## Definition (Johnstone construction/plus construction)

Let  $\mathcal{F}, \mathcal{G} \in \text{Sh}(X)$  and  $f : \mathcal{F} \rightarrow \mathcal{G}$ , then

- $\mathcal{F}^+ = \{S \subseteq \mathcal{F} \mid \square(\ulcorner S \text{ is singleton} \urcorner)\} / \sim$ , where  $S \sim T \Leftrightarrow \square(S = T)$
- $f^+ : \mathcal{F}^+ \rightarrow \mathcal{G}^+$ ,  $[S] \mapsto [\{f(x) \mid x \in S\}]$
- $\gamma : \mathcal{F} \rightarrow \mathcal{F}^+$ ,  $x \mapsto [\{x\}]$

## Proposition

For a sheaf  $\mathcal{F}$ ,  $\mathcal{F}^{++}$  is a  $\square$ -sheaf.

# How to construct $\varphi^{++}$

## Proposition

Let  $\square$  be a modal operator,  $\varphi$  be some formula, Then  $\varphi^{\square} \Leftrightarrow ((\varphi^{\square})^+)^+$  intuitionistically.

## Definition (+-translation)

- We replace  $\bar{x} \in \mathcal{F}$  by  $\bar{x} \in \mathcal{F}^+$ .
- Terms  $x_1, \dots, x_n$  and  $f(\bar{y})$  are replaced by  $\gamma(x_1), \dots, \gamma(x_n)$  and  $f^+(\bar{y})$ .
- $\varphi^+$  of a formula  $\varphi : \mathcal{F} \rightarrow \Omega$  is attained by replacing all free variables with their  $\gamma$ -images and morphisms and domains of quantifications with their +-constructions, e.g.  
 $(\forall x : \mathcal{F} f(x) = g(x))^+ := \forall x : \mathcal{F}^+ f^+(x) = g^+(x)$ .

# Conclusion

## Theorem

Let  $X$  be a topological space and  $\square$  be a modal operator on  $Sh(X)$  then for any formula  $\varphi$

$$X \models_{Sh(X)} (\varphi^\square)^{++} \iff X_\square \models_{Sh(X_\square)} \varphi^{++}$$

Now when proving the theorem we can pretend that the context and that all domains of quantification are  $\square$ -sheaves.

## Theorem

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